## eFLINT - A DSL for Testing Normative Specifications

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# Policy-CAD-----



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Deontic	Potestative				
duties, obligations	powers, actions				
permissions	liabilities				

# Analyzing legal cases



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How do we write down an interpretation formally?





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## Formal Language for the Interpretation of Norms (FLINT)

Robert van Doesburg / Tijs van der Storm / eFLINT

#### Commonalities

- Judgements characterize the relevant sub-set of the world
  - certain facts are postulated (to hold true or false)
  - other facts are derived (from other judgements)
- Transition rules determine the availability of actions and their effects

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#### Challenges

- Language design: appeal, scope, fit-for-purpose ...
- Policy design: consistency, composition, qualification ...
- Policy analysis: exploration, testing, verification, reasoning, planning ...
- System compliance: testing, verification, "by construction" ...

- World: values, types, expressions
- Orms: duties, acts, transitions
- O Pragmatics: refinement, scripts, testing

Fact-type declarations associate a type with a fact identifier:

1	Fact	citizen			
2	Fact	candidate	ldentified	by	Atom
3	Fact	administrator	ldentified	by	Atom
4	Fact	voter	ldentified	by	citizen
5	Fact	winner	ldentified	by	candidate
6	Fact	vote	ldentified	bу	(voter * candidate)

Types are essentially record-types:

• Field names are variables (possibly decorated fact identifiers)

```
1 Alice
2 7
3 4 Alice:citizen
5 Chloe:candidate
6 Admin:administrator
7 8 (Alice:citizen):voter
9 10 ((Alice:citizen):voter, Chloe:candidate):vote
```

example instances

```
1 Alice
2 7
3 4 Alice:citizen
5 Chloe:candidate
6 Admin:administrator
7 8 (Alice:citizen):voter
9 10 ((Alice:citizen):voter, Chloe:candidate):vote
```

example instances

The state of the world at any particular moment is a set of instances  $\sigma$ , containing those instances that *hold true* at that moment

#### Expressions

• Expressions evaluate to atoms, integers, Booleans or instances of fact-types

```
citizen
2
3
   citizen (Alice)
4
   voter(citizen(Alice))
5
   voter (Alice)
6
7
   voter(citizen = citizen(Alice))
8
   vote(voter(Alice), Chloe)
   vote (voter = voter (Alice), candidate = Chloe)
9
10
   vote(candidate = Chloe, voter = voter(Alice))
11
12
   vote(voter = voter(Alice))
13
   vote(candidate = candidate, voter = voter(Alice))
14
   vote()
```

variables and constructors

```
1 Holds(voter(Alice))
2
3 vote[voter]
4 vote[candidate]
5
6 vote[candidate] When Holds(vote)
7 vote[candidate] When vote
```

operators

1 2 3 Quantifiers bind variables to all instances of the variable's type:

```
(Exists candidate : vote(voter(Alice), candidate))
```

```
(Forall citizen : vote(voter(citizen), Chloe))
```

Foreach can only be used in combination with an *aggregator*.

1 Count(Foreach vote : vote When Holds(vote) && vote[candidate] = candidate)

1 2 Derivation expression as a predicate (type-components are bound):

```
Fact has voted Identified by voter
```

```
Holds when (Exists candidate : vote(voter, candidate))
```

1	Predicate	vote concluded When (Exists candidate : winner(candidate))	
2	Predicate	voters done When (Forall citizen : !voter()    has voted(v	oter()))

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Derivation expression computes the set of instances that hold true:

- Fact number of votes Identified by Int
  - Derived from Count(Foreach vote : vote When Holds(vote))

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Derivation expression computes the set of instances that hold true:

```
Fact number of votes Identified by Int
```

```
Derived from Count(Foreach vote : vote When Holds(vote))
```

```
• Derived facts cannot be postulated
```

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#### Recall Hohfeld's conceptions



fundamental relation: **duty-claim** between *duty holder* and *claimant* 

fundamental relation: **power-liability** between *actor* and *recipient* 

How do we write down an interpretation formally?

#### A duty indicate that its holder ought to perform some action:

- 1 Duty cast vote duty Holder voter Claimant administrator
  - A duty-type declaration is a fact-type declaration with a record-type

Actions modify the world by adding or removing instances from  $\sigma$ :

```
1 Act cast vote
2 Actor voter
3 Recipient administrator
4 Related to candidate
5 Conditioned by voter && !has voted()
6 Creates vote()
7 Terminates cast vote duty()
```

• An act-type declaration is a fact-type declaration with a record-type

A transition is a state  $\sigma$ , an instance *a* of an act, and the sets *T* and *C* of instances terminated and created by the act, if and only if:

- $\textcircled{O} a \text{ holds true in } \sigma$
- 2 the pre-condition of a holds in  $\sigma$
- **③** T is the result of evaluating the terminating post-conditions of a in  $\sigma$
- **③** C is the result of evaluating the creating post-conditions of a in  $\sigma$

$$(\sigma) \longrightarrow \langle a, T, C \rangle$$

Derived facts may have to be recomputed after an action is performed:

```
1 Act cast vote
2 Actor voter
3 Recipient administrator
4 Related to candidate
5 Conditioned by voter && !has voted()
6 Creates vote()
7 Terminates cast vote duty()
```

•  $\sigma'$  may be incomplete and inconsistent w.r.t. derivation expressions

$$(\sigma) \xrightarrow[\langle a, T, C \rangle]{\sigma'} \xrightarrow[]{\sigma'} \xrightarrow[]{\sigma'} \xrightarrow[]{\sigma'}$$

1

```
Act enable vote
2
3
     Actor administrator
     Recipient citizen
4
     Conditioned by !voter() && !vote concluded()
5
     Creates voter(),
6
7
             cast vote duty(voter = voter()),
             (Foreach candidate : cast vote(voter = voter()))
```

# Completing the example (2)

Placeholders can be introduced for types:

1 Placeholder other candidate For candidate

```
Act declare winner
2
     Actor administrator
3
     Recipient candidate
     Conditioned by
4
5
          !vote concluded()
6
7
      && voters done()
      && (Forall other candidate :
8
             Count(Foreach vote : vote[voter]
9
                     When vote && vote [candidate] == other candidate) <
10
             Count(Foreach vote : vote[voter]
11
                     When vote && vote [candidate] == candidate)
12
           When other candidate != candidate)
13
     Creates winner(candidate)
```

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A refinement of a policy description replaces all simple, infinite types with finite types:

1	Fact	citizen	ldentified	by	[John,	Frank,	Peter,	Chloe ,	Hannah]
2	Fact	candidate	ldentified	by	[Mary,	David]			
3	Fact	administrator	ldentified	by	Admin				

#### and also identifies an initial state (implicit Foreach):

```
    administrator.
    citizen.
    candidate.
    declare winner(Admin, candidate).
    enable vote(Admin, citizen).
```

#### • A refinement enables exploring the reachable states manually

• Scripts are basic programs for stepping through reachability graphs

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Action call !<EXPR>. evaluates to an enabled act and executes it (or fails)

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Action call !<EXPR>. evaluates to an enabled act and executes it (or fails)

Query ?<EXPR>. fails if expression does not evaluate to True in the current state

```
!enable vote(citizen = John).
2
  !enable vote(citizen = Frank).
3
  !enable vote(citizen = Peter).
4
  !cast vote(voter = voter(John), candidate = Mary).
5
  !cast vote(voter = voter(Frank), candidate = Mary).
  !cast vote(voter = voter(Peter), candidate = David).
6
7
  !declare winner().
8
  ?winner(Mary).
9
  ?(Forall candidate : !winner() When candidate != Mary).
```

```
1 !enable vote(citizen = John).
2 !enable vote(citizen = Frank).
3 !enable vote(citizen = Peter).
4 !cast vote(voter = voter(Frank), candidate = Mary).
5 !cast vote(voter = voter(Peter), candidate = David).
6 !enable vote(citizen = Hannah).
7 !cast vote(voter = voter(Hannah), candidate = David).
8 !declare winner().
```

```
!enable vote(citizen = John).
2
  !enable vote(citizen = Frank).
3
  !enable vote(citizen = Peter).
  !cast vote(voter = voter(Frank), candidate = Mary).
4
5
  !cast vote(voter = voter(Peter), candidate = David).
  !enable vote(citizen = Hannah).
6
  !cast vote(voter = voter(Hannah), candidate = David).
7
8
  !cast vote(voter = voter(John), candidate = Mary).
9
  !declare winner().
```

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Curb your enthusiasm Thomas...

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```
public void register voter(String name) {
  registered voters.add(name);
  ScriptGen.add actor(name, "citizen"); // gualification
  ScriptGen.enable vote(name);
public void register candidate(String name) {
  vote count.put(name.0):
  ScriptGen.add actor(name, "candidate"); // qualification
public void vote(String voter. String candidate) {
  vote count.put(candidate, vote count.get0rDefault(candidate.0) + 1):
  ScriptGen.cast vote(voter, candidate); // gualification
public String find winner() { // returns null if there is no winner
  String winner = null:
  int winning votes = 0:
  for (Entry<String, Integer> entry : vote count.entrySet()) {
    if(entry.getValue() > winning votes) {
     winner = entrv.getKev():
      winning votes = entry.getValue();
  if (winner != null) { // gualification
    ScriptGen.declare winner(winner):
  return winner:
```

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